# Hyper-Threading Technology and its Impact on OpenMP\*



## Agenda

- Hyper-Threading Overview
  - Exploiting Hyper-Threading Technology
    - Explicit Threads
    - OpenMP Programming API
  - OpenMP\* Programming Example



## Today's Processors

- Single Processor Systems
  - Instruction Level Parallelism (ILP)
  - Performance improved with more CPU resources

Hyper-Threading technology enables TLP to single processor system.

- Multiprocessor Systems
  - Thread Level Parallelism (TLP)
  - Performance improved by adding more CPUs

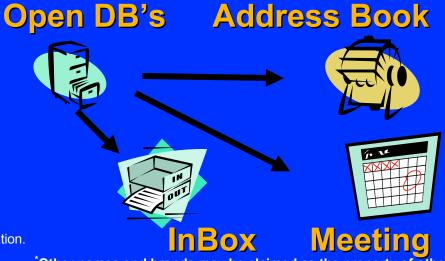


## Today's Software

Sequential tasks



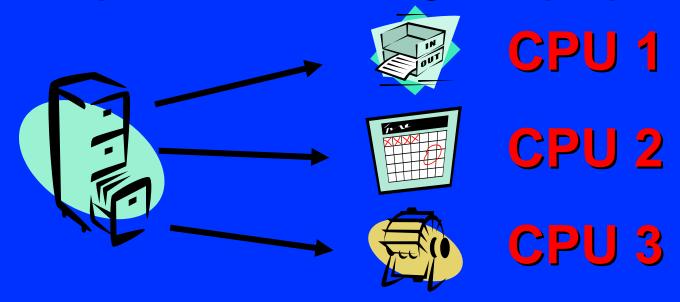
Parallel tasks





## Multi-Processing

Run parallel tasks using multiple processors



Multi-tasking workload + processor resources => Improves MT Performance



# Hyper-Threading Overview The Increase in Instruction Processing (Multi-tasking Workload from Previous Slide) Throughput of Hyper-Threading is Due to:

- The design of the Intel Netburst Microarchitecture
- The mix of IA-32 Instructions typically found in multi-threaded code



## Why Hyper-Threading?

- Studies[1] show that only 35% of execution resources of the Intel NetBurst Micro-architecture are used
- Hyper-Threading technology takes advantage of the inherent parallelism of multithreaded code to provide the processor core with a second thread of execution



### Hyper-Threading Technology

Multiprocessor

**Hyper-Threading** 

Processor
Execution
Resources

Processor
Execution
Resources

Processor
Execution
Resources

AS = Architecture State (eax, ebx, control registers, etc.), xAPIC

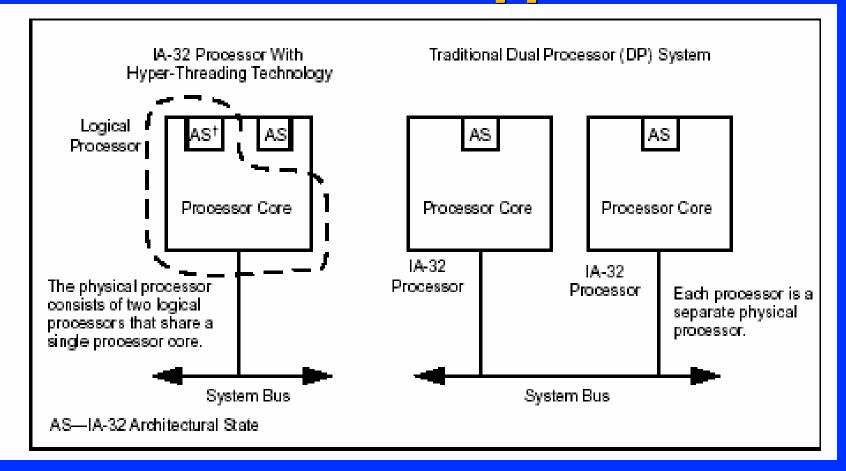
Hyper-Threading Technology looks like two processors to software

### How is this done?

- Processors that are enabled with Hyper-Threading Technology:
  - –Manage incoming instructions from two different software threads
  - The processor keeps track of the data processing status of each set of instructions



## Switching from One Logical Processor to the Other [1]



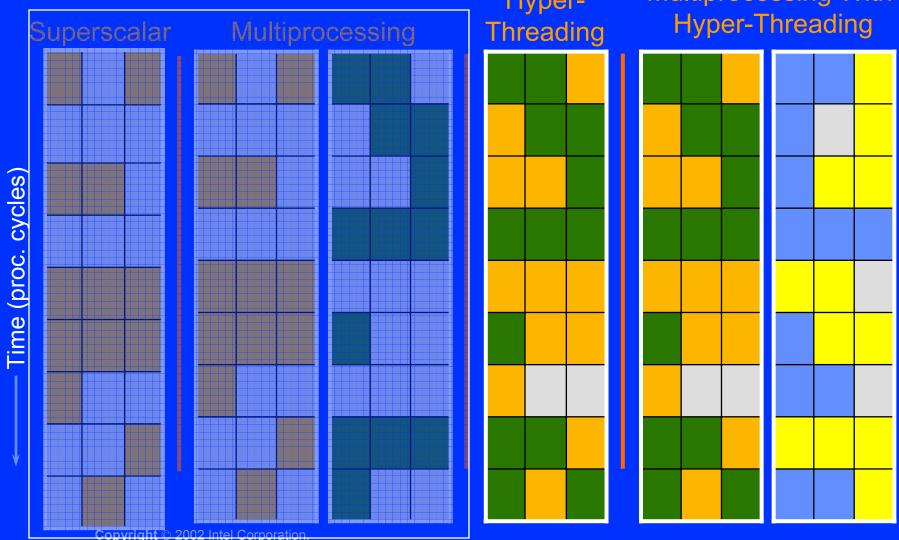


## Three Intel Xeon<sup>TM</sup> Design Goals to Support Hyper-Threading[2]

- Minimize Die Area less than 5% in additional die area cost
- When one logical processor stalls the other logical processor continues to make forward progress
- A single threaded application running on a processor with Hyper-Threading technology executes at same speed as a processor without this capability



Resource Utilization Multiprocessing With Hyper- Threading



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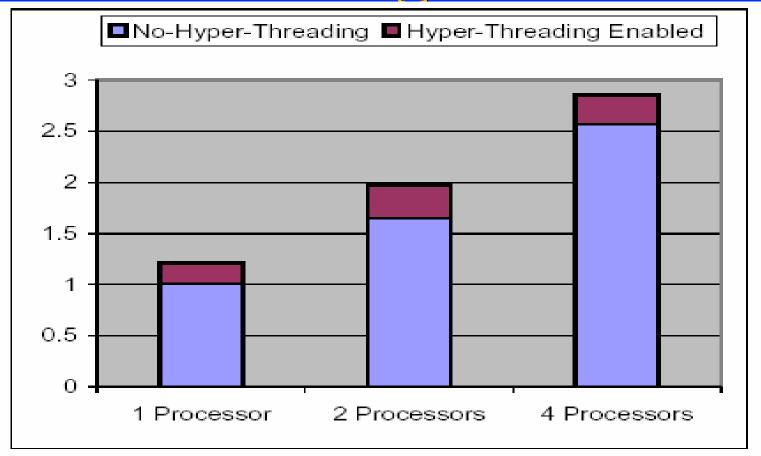
Note: Each box represents a processor execution unit Other names and brands may be claimed as the property of others.

## Key Point

- Hyper-Threading technology enables better utilization of hardware resources
- Hyper-Threading technology provides more computing power for multithreaded applications

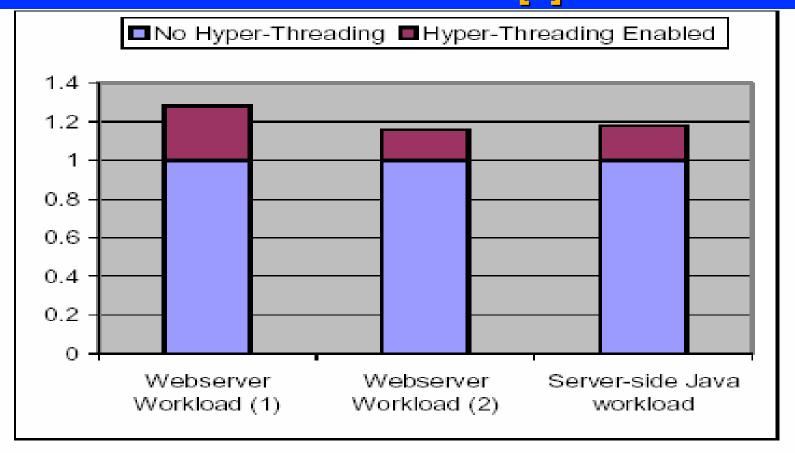


## Performance Increases from Hyper-Threading Technology on an Online Transaction Processing Workload[2]





## Web Server Benchmark Performance[2]





# Hyper-Threading Overview Software-based Speculative Precomputation[3]

- Technique to improve the latency of single-threaded applications
- Algorithmic Sketch:
  - Speculative thread fetches memory objects (usually in a strided manner)
  - Main thread does the computation with prefetched data objects
- Acronym SP



#### Initial performance data: speculative prefetching (SP) on a preproduction version of an Intel® Xeon<sup>TM</sup> processor with Hyper-Threading Technology[3]

Benchmark	Description	Speed-up
Synthetic	Graph traversal in large random graph simulating large database retrieval	22% - 45%
MST (Olden)	Minimal Spanning Tree algorithm used for data clustering	23% - 40%
Health (Olden)	Hierarchical database modeling health care system	11% - 24%
MCF (SPEC2000int)	Integer programming algorithm used for bus scheduling	7.08 %



## Hyper-Threading Overview Performance Gains with Hyper-Threading Technology

- Hyper-Threading technology can provide a performance gain of up to 30% over a comparable IA-32 processor without Hyper-Threading technology, assuming:
  - Multithreaded operating system and application code
- For multiprocessor systems:
  - Increase in computing power will generally scale linearly with an increase in the number of physical processors
- Scalability of performance is highly dependent on the nature of the application



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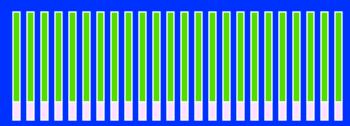


#### **Parallel Computing Overview**

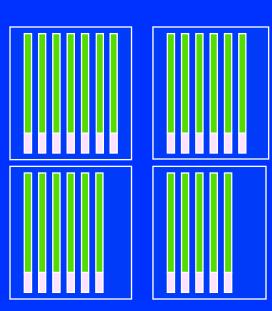
#### Writing a Parallel Application



Decompose into tasks



Original Problem



Group onto units.

Tasks, shared and local data

Code with a parallel Prog. API

Units of execution + new shared data for extracted dependencies

be cla Corresponding source code

## Parallel Computing Overview What is OpenMP\*?

C\$OMP FLUSH

#pragma omp critical

C\$OMP THREADPRIVATE (/ABC/)

CALL OMP SET NUM THREADS (10)

## OpenMP\*: An API for Writing Multithreaded Applications

- Compiler directives and library routines for parallel application programmers
- Makes it easy to create multi-threaded (MT) programs in Fortran, C and C++
- Standardizes last 15 years of SMP practice

Com Inched Collin () Dir. / )

CŞOMP DO lastprivate(XX)

Nthrds = OMP\_GET\_NUM\_PROCS()

omp set lock(lck)

intط

\*Other names and bra.....rs

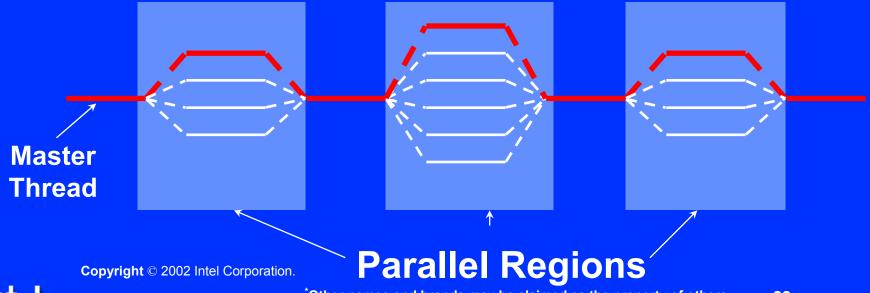
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#### **Parallel Computing Overview**

### OpenMP\* Programming Model

#### Fork-Join Parallelism:

- Master thread spawns a team of threads as needed.
- Parallelism is added incrementally: i.e., the sequential program evolves into a parallel program.





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## Pi Program

```
static long num_steps = 100000;
double step;
void main ()
       int i; double x, pi, sum = 0.0;
       step = 1.0/(double) num steps;
       for (i=1;i<= num steps; i++){
              x = (i-0.5)*step;
              sum = sum + 4.0/(1.0+x*x);
       pi = step * sum;
```



### Pi: Windows\*

## Tanceacs

```
#define NUM THREADS 2
HANDLE thread handles[NUM THREADS];
CRITICAL SECTION hUpdateMutex;
double global sum = 0.0;
void Pi (void *arg)
  int i, start;
 double x, sum = 0.0;
 static long num steps = 100000;
 double step;
 start = *(int *) arg;
 step = 1.0/(double) num steps;
 for (i=start;i<= num steps; i=i+NUM THREADS){
    x = (i-0.5)*step;
    sum = sum + 4.0/(1.0+x*x);
 EnterCriticalSection(&hUpdateMutex);
 global sum += sum;
 LeaveCriticalSection(&hUpdateMutex);
```

```
void main ()
 double pi; int i;
 DWORD threadID:
 int threadArg[NUM THREADS];
 for(i=0; i<NUM THREADS; i++) threadArg[i] = i+1;
 InitializeCriticalSection(&hUpdateMutex);
 for (i=0; i<NUM THREADS; i++){
        thread handles[i] = CreateThread(0, 0,
                         (LPTHREAD START ROUTINE) Pi,
                         &threadArg[i], 0, &threadID);
 WaitForMultipleObjects(NUM THREADS,
                         thread handles, TRUE, INFINITE);
 pi = global sum * step;
 printf(" pi is %f \n",pi);
```

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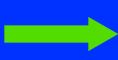


## Simple Is Better

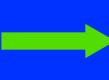
#### **Threads libraries:**

- Pro: Programmer has control over everything
- Con: Programmer <u>must</u> control everything

Control over all threads



High complexity



High programming costs

The simplicity of OpenMP\* lowers programming costs.



## Pi: OpenMP\* version

```
#include <omp.h>
 static long num steps = 100000;
                                       double step;
 void main ()
         int i; double x, pi, sum = 0.0;
          step = 1.0/(double) num steps;
#pragma omp parallel for reduction(+:sum) private(x)
          for (i=1;i \le num \text{ steps}; i++)
                 x = (i-0.5)*step;
                 sum = sum + 4.0/(1.0 + x*x);
          pi = step * sum;
                                    OpenMP* adds 2
                                       lines of code.
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```



## OpenMP\*: Easy as Pi

```
#include <omp.h>
static long num steps = 100000;
                                     double step;
void main ()
        int i; double x, pi, sum = 0.0;
        step = 1.0/(double) num steps;
#pragma omp parallel for reduction(+:sum) private(x)
        for (i=1;i \le num \text{ steps}; i++)
                x = (i-0.5)*step;
                sum = sum + 4.0/(1.0+x*x);
        pi = step * sum;
```

OpenMP\* simplifies multithreading.



## Key Take Away

If a multi-threaded application performance does not improve on an MP system, you will have minimal benefit with Hyper-Threading technology.



## Summary

- Hyper-Threading Technology gives you more computing power to throw at your problems.
- OpenMP\* is an easy to use API for writing multithreaded programs.
- Continue to use good threaded programming practices with Hyper-Threading technology.



### References

- [1] "Introduction to Hyper-Threading Technology" at the URL: http://www.intel.com/technology/hyperthread/download/25000802.pdf.
- [2] D. T. Marr, F Binns, D. L. Hill, G. Hinton, D. Koufaty, J. A. Miller, M. Upton, "Hyper-Threading Technology Architecture and Microarchitecture", *Intel Technology Journal*, Vol. 6, No. 1, February 2002.
- [3] H. Wang, P. Wang, R. D. Weldon, S. M. Ettinger, H. Saito, M. Girkar, S. S. Liao, J. P. Shen, "Speculative Precomputation: Exploring the Use of Multithreading for Latency", *Intel Technology Journal*, Vol. 6, No. 1, February 2002.



### Call to Action

- Think of Hyper-Threading as a Technology to Improve Instruction Throughput of Processors
- Can Benefit Multiprocessor Applications, e.g. OpenMP\*

